



Linux/ia64 support for performance monitoring

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Outline

- Goals of the interface
- The perfmon-2 kernel interface overview
- Overview of existing tools
- Examples of measurements

Goals of the interface

- Provides a **generic** interface to access PMU
 - Not dedicated to one app, avoid fragmentation
- Must be portable across PMU models:
 - Almost all PMU-specific knowledge in user level libraries
- Supports **per-process** monitoring
 - Self-monitoring, unmodified binaries, attach/detach
 - multi-threaded and multi-process workloads
- Supports system-wide monitoring
- Supports counting and sampling
- No modification to applications or system
- **Builtin**, efficient, robust, secure, simple, documented

Current implementations

- In Linux/ia64 since 2.4.0
- In all 2.4-based kernels: perfmon-1
 - First generation interface
 - Included in SLES-8, RHAS-2.1, RHEL-3.0 (but broken)
 - Several limitations : no monitoring across fork()
- In all 2.6-based kernels: perfmon-2
 - Second generation interface
 - **Not backward compatible** with perfmon-1
 - Lots of new features and more flexibility
 - Easier to port existing applications (Oprofile, VTUNE)

Perfmon-2 new features

- Perfmon context identified with file descriptor
- Monitoring across pthread_create/fork (i.e., clone2)
 - Using ptrace() to help track creation/exit
- Custom sampling buffer formats
 - Easy to port existing monitoring tools
- Can attach/detach to/from running process
 - Useful for long running daemons
- Event-set multiplexing
 - Minimize effect of limited number of counters
- Scalable counter overflow notification mechanism
 - Exploit Linux file system interface

Perfmon-2 interface

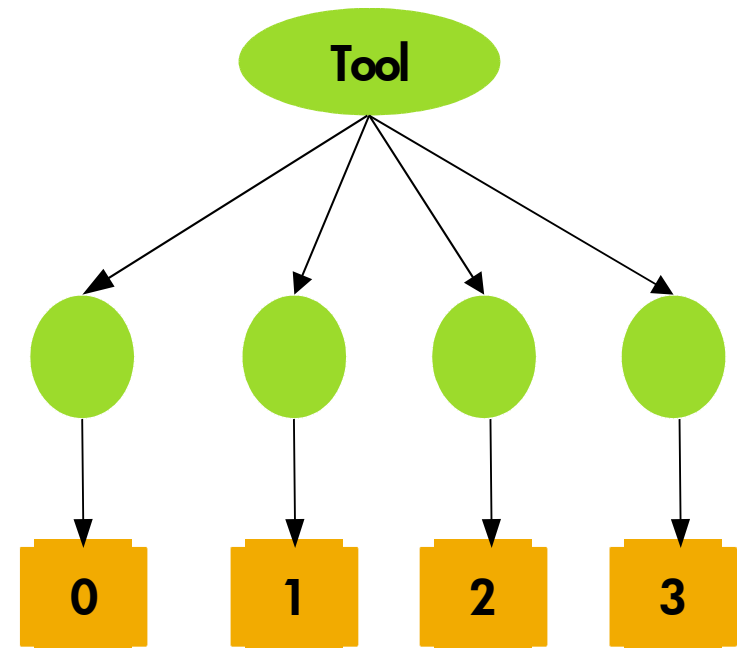
- Uses a **system call** instead of drivers interface
 - More flexibility (e.g., # arguments)
 - Close ties with the context switch code, exit, fork
 - Kernel compile-time option (CONFIG_PERFMON)
- Perfmon **context** to encapsulate PMU state
 - Each context uniquely identified by file descriptor
- Argument similar to ioctl() but you can pass vectors
- Counters are always exported as 64-bit wide

int perfmonctl(int fd, int cmd, void *arg, int nargs)

PFM_CREATE_CONTEXT	PFM_READ_PMDS	PFM_START
PFM_WRITE_PMCS	PFM_LOAD_CONTEXT	PFM_STOP
PFM_WRITE_PMDS	PFM_UNLOAD_CONTEXT	PFM_RESTART

System wide monitoring

- Monitoring of all processes
- Union of cpu-wide sessions:
 - flexibility: different events on each CPU
 - simplicity for implementations
 - Uses process pinning capabilities
- Ability to exclude the idle task
- Does not currently coexist with per-process sessions



Support for event-based sampling

- 64-bit counter overflow notification via message
 - Extracted via read(), supports select/poll, SIGIO
- Number of sampling periods = number of counters
 - Sampling period can be randomized (reproduceable)
 - Each counter has long and short period to hide recovery cost
- task can be stopped on overflow notification
- Kernel level sampling buffer
 - Variable size
 - Remapped to user level to avoid massive copying
- Sampling buffer format controlled by kernel module
 - Perfmon core does not know how samples are recorded

Custom sampling buffer formats

- Why?
 - No single format can satisfy all needs
 - Make it easy to port existing monitoring tools
- How?
 - Each format is implemented as a kernel module
 - Each format is uniquely identified with 128-bit UUID
 - Each format registers a set of callbacks with perfmon core
 - Each format provides a handler callback invoked on counter overflow
- What can a format use?
 - May use buffer allocation+remapping service
 - May use overflow notification mechanism

Event set multiplexing

- To overcome limitation on number of counters
- Each set encapsulates PMU state (PMC/PMD/IBR/DBR)
 - Can define up to 65K sets
 - Possiblity to exclude idle task per set (system-wide only)
- Supported in per-task and system-wide modes
 - counting and sampling support
- Cycling and Switching:
 - Cycling is round-robin
 - Per-set time-based: timeout granularity of clock tick
 - Per-set overflow-based: “after n overflows of counter X”
 - Multiple counters can be triggers in overflow mode
- Can be used to implement counter cascading

Linux/ia64 monitoring tools

- gprof:
 - Flat profile, call graph, recompilation, applications only
- Kernel profiler:
 - Kernel only flat profile, builtin (cmdline profile=)
- Prospect(HP)/OProfile:
 - PMU-based, system-wide flat profile
- VTUNE(Intel):
 - PMU-based, system-wide flat profile, Windows-side GUI
- pfmon/libpfm, qprof, q-syscollect/q-view (HP Labs)
- PAPI toolkit (U. of Tennessee):
 - PMU-based, counting, sampling, uses libpfm

Monitoring complicated workloads

- Can follow across fork/vfork, pthread_create
- Works for counting and sampling
- Supports regular expression to filter binaries of interest
- Example: elapsed cycles for a simple compile

```
$ pfmon --us-c -u -k --follow-all -ecpu_cycles,ia64_inst_retired \  
-- cc e.c -o e
```

1,164,772	CPU_CYCLES	/usr/lib/gcc-lib/ia64-linux/2.96/cpp0
1,295,480	IA64_INST_RETIRED	/usr/lib/gcc-lib/ia64-linux/2.96/cpp0
13,758,346	CPU_CYCLES	/usr/lib/gcc-lib/ia64-linux/2.96/cc1
21,863,635	IA64_INST_RETIRED	/usr/lib/gcc-lib/ia64-linux/2.96/cc1
5,708,731	CPU_CYCLES	as
7,165,599	IA64_INST_RETIRED	as
27,046,535	CPU_CYCLES	/usr/bin/ld
35,247,760	IA64_INST_RETIRED	/usr/bin/ld
1,381,134	CPU_CYCLES	/usr/lib/gcc-lib/ia64-linux/2.96/collect2
1,508,977	IA64_INST_RETIRED	/usr/lib/gcc-lib/ia64-linux/2.96/collect2
1,913,253	CPU_CYCLES	cc
1,976,590	IA64_INST_RETIRED	cc

Detailed cycle breakdown

- Can use current pfmon with wrapper script
 - i2prof.pl written by Per Ekman
- Using the current development version of pfmon:
 - exploits event sets multiplexing capabilities

```
$ pfmon -m itanium2-stalls -k -u -system-wide -print-interval - mcf inp.in
```

#	-----											
#	%itlb	%icache	%bra	%unstall	%BE	%score	%RSE	----- D-access -----				
#				exec	flush	board		%d1tlb	%d2tlb	%cache	-loaduse-	
#										res	%gr	%fr
#	-----											
	0.00	0.02	2.81	32.08	10.06	1.19	0.00	0.57	5.28	4.19	43.80	0.00
	0.00	0.02	2.81	32.12	10.06	1.19	0.00	0.57	5.28	4.19	43.77	0.00
	0.00	0.02	2.81	32.09	10.06	1.19	0.00	0.57	5.28	4.19	43.78	0.00
	0.00	0.00	0.08	59.29	0.22	0.05	0.00	0.03	0.01	1.75	38.57	0.01
	0.00	0.00	0.06	54.49	0.16	1.16	0.00	0.46	3.16	3.74	36.76	0.00
	0.00	0.05	2.83	42.14	10.08	1.06	0.02	0.68	4.77	5.69	32.69	0.00
	0.00	0.05	2.79	42.27	9.97	1.07	0.02	0.69	4.88	5.67	32.59	0.00
	0.00	0.03	2.44	41.42	8.74	1.11	0.00	0.55	4.30	4.32	37.09	0.00
	0.00	0.02	2.82	32.07	10.07	1.16	0.00	0.62	5.69	4.46	43.08	0.00

Data load cache misses profiles

- Obtained using the Data EARS
- Provides instruction AND data views
- Careful with interpretation: not all misses lead to stalls
- To be included in the next version of pfmon
- Example: mcf instruction and data views

#count	%self	%cum	%L2	%L3	%RAM	instruction addr
6358	11.11%	11.11%	3.05%	5.17%	91.77%	price_out_impl+0x820<mcf>
6238	10.90%	22.01%	26.74%	69.93%	3.33%	price_out_impl+0x850<mcf>
5404	9.44%	31.45%	74.43%	24.94%	0.63%	bea_compute_red_cost+0x50<mcf>
5016	8.77%	40.22%	46.69%	33.77%	19.54%	bea_compute_red_cost+0xa1<mcf>
4968	8.68%	48.90%	42.43%	9.98%	47.58%	primal_bea_mpp+0x7b1<mcf>
4878	8.52%	57.42%	36.67%	51.87%	11.46%	bea_compute_red_cost+0x90<mcf>

#count	%self	%cum	%L2	%L3	%RAM	data addr
37	0.06%	0.06%	62.16%	32.43%	5.41%	0x2000000000017ebd0
32	0.06%	0.12%	75.00%	18.75%	6.25%	0x200000000000d07b0
29	0.05%	0.17%	68.97%	24.14%	6.90%	0x200000000000e2438
28	0.05%	0.22%	96.43%	3.57%	0.00%	0x200000000000d3708
26	0.05%	0.27%	88.46%	11.54%	0.00%	0x200000000000d8c58

Kernel level call stack sampling

- Combines kernel stack unwinder with perfmon:
 - On counter overflow, record the call stack
 - Uses a custom sampling buffer format
- Example using the development version of pfmon:

```
$ pfmon -el3_misses --long-smpl-periods=2000 --smpl-periods-random=0xff:10 -k \  
--smpl-module=kcall-stack-ia64 --resolve-addr --system-wide
```

```
__copy_user, file_read_actor, do_generic_mapping_read, __generic_file_aio_read, generic_file_aio_read,  
do_sync_read, vfs_read, sys_read, ia64_ret_from_syscall
```

```
do_anonymous_page, do_no_page, handle_mm_fault, ia64_do_page_fault, ia64_leave_kernel
```

```
clear_page, do_anonymous_page, do_no_page, handle_mm_fault, ia64_do_page_fault, ia64_leave_kernel
```

```
bh_lru_install, __find_get_block, __getblk, ext3_get_inode_loc, ext3_reserve_inode_write,  
ext3_mark_inode_dirty, ext3_dirty_inode, __mark_inode_dirty, update_atime, link_path_walk, open_namei,  
filp_open, sys_open, ia64_ret_from_syscall
```

```
end_bio_bh_io_sync, bio_endio, __end_that_request_first, scsi_end_request, scsi_io_completion,  
sd_rw_intr, scsi_finish_command, scsi_softirq, do_softirq, ia64_handle_irq, ia64_leave_kernel
```

```
filemap_nopage, do_no_page, handle_mm_fault, ia64_do_page_fault, ia64_leave_kernel
```

```
scsi_finish_command, scsi_softirq, do_softirq, ia64_handle_irq, ia64_leave_kernel
```

```
end_page_writeback, end_buffer_async_write, end_bio_bh_io_sync, bio_endio, __end_that_request_first,  
scsi_end_request, scsi_io_completion, sd_rw_intr, scsi_finish_command, scsi_softirq, do_softirq,  
ia64_handle_irq, ia64_leave_kernel
```

Data structure profiling(experimental)

- Statistical profile of data structure to improve layout
- combines loads+stores retired events, range restrictions
- Currently user level but can be made to work in kernel

```
static void
walk_list(node_t *list)
{
    while(list) {
        list->total = list->val1 + list->val2;
        list = list->next;
    }
}
```

field name	:size:	offs	: line	: cover	: loads	: stores	:
node_t->prev	: 8 :	8	: 0	: 19.98%	: 0.00%	: 0.00%	:
node_t->next	: 8 :	0	: 0	: 19.67%	: 32.86%	: 0.00%	:
node_t->val1	: 8 :	16	: 0	: 20.35%	: 33.97%	: 0.00%	:
node_t->val2	: 8 :	24	: 0	: 20.00%	: 33.17%	: 0.00%	:
node_t->total	: 8 :	32	: 0	: 20.01%	: 0.00%	: 100.00%	:

Conclusions

- The Itanium® 2 PMU is very powerful
- Monitoring is key to achieving world-class performance
- Linux/ia64 has the most advanced monitoring subsystem of all Linux implementations
- Linux/ia64 already supports a variety of monitoring tools
- We need to develop better, smarter tools for non-experts

Linux/ia64
perfmon resources

Linux/ia64 perfmon resources

- Performance tools developed by HPLabs:
 - pfmon/libpfm, q-tools, q-prof
 - <http://www.hpl.hp.com/research/linux>
- Intel VTUNE:
<http://www.intel.com/software/products/vtune>
- PAPI
<http://icl.cs.utk.edu/projects/papi>
- Prospect:
<http://prospect.sf.net>
- OProfile
<http://oprofile.sf.net>

Linux/ia64 perfmon resources

- i2prof.pl:
<http://www.pdc.kth.se/~pek/i2prof.pl>
- IPF PMU architecture:
<http://developer.intel.com/design/itanium/>
- Itanium® 2 PMU specification:
<http://developer.intel.com/design/itanium/manuals.htm>

Backup slides

Opcode matching with pfmon

- Constrains monitoring to instructions or patterns
 - Based on opcode, e.g., st8.*
 - Based on functional unit, e.g., M,F,I,B
 - Pattern uses a match+mask fields
 - Not all instructions can be uniquely identified
 - Two opcode matching registers on Itanium® 1 & 2
- Ex.: counting the number of br.cloop instructions:

```
$ pfmon -us-c --opc-match8=0x1400028003ffff1fa \  
  -e IA64_TAGGED_INST_RETIRED_IBRP0_PMC8 -- foo  
      4,999,950,164 IA64_TAGGED_INST_RETIRED_IBRP0_PMC8
```

Range restrictions

- Constrains monitoring to range of data or code
 - Implemented via debug registers (not used as breakpoints)
 - Can specify a range inside the kernel
 - Works for both per-process and system-wide
 - Not all events support range restrictions
- Range must be aligned on size for exact measurements
 - gcc -falign-functions= option can be useful
- Pfmmon supports symbol name for the range
- Ex.: how many L2 misses while executing init_tab()

```
$ pfmmon -us-c -el2_misses -- foo
      1,245,516 L2_MISSES (misses for the entire execution)

$ pfmmon -us-c -irange=init_tab -el2_misses -- foo
      14,456 L2_MISSES (misses for init_tab() only)
```

Sampling branches (BTB)

- Capture up to the last 4 branches:
 - Each entry contains source/target addr., prediction outcome
 - Possible to filter branches: taken/not taken, mispredicted
 - Can be combined with EAR to build a path to a cache/tlb miss
- Ex.: sample every 1000 taken branch, record last 4

```
$ pfmon --smpl-periods-random=5:0xff --btb-tm-tk \  
      --long-smpl-periods=1000 -ebranch_event -- foo
```

```
entry 231 PID:673 CPU:0 STAMP:0x12957325ac49 IIP:0x400000000000004d0  
  last reset : 1004  
  branch source address: 0x400000000000004f2  
  branch target address: 0x400000000000004c0  
  branch taken : yes, prediction: success, pipe flush: no  
  ...  
400000000000004f0: [MFB]          nop.m 0x0  
400000000000004f1:          nop.f 0x0  
400000000000004f2:          br.cloop.sptk.few 400000000000004c0
```

Simple self-monitoring example

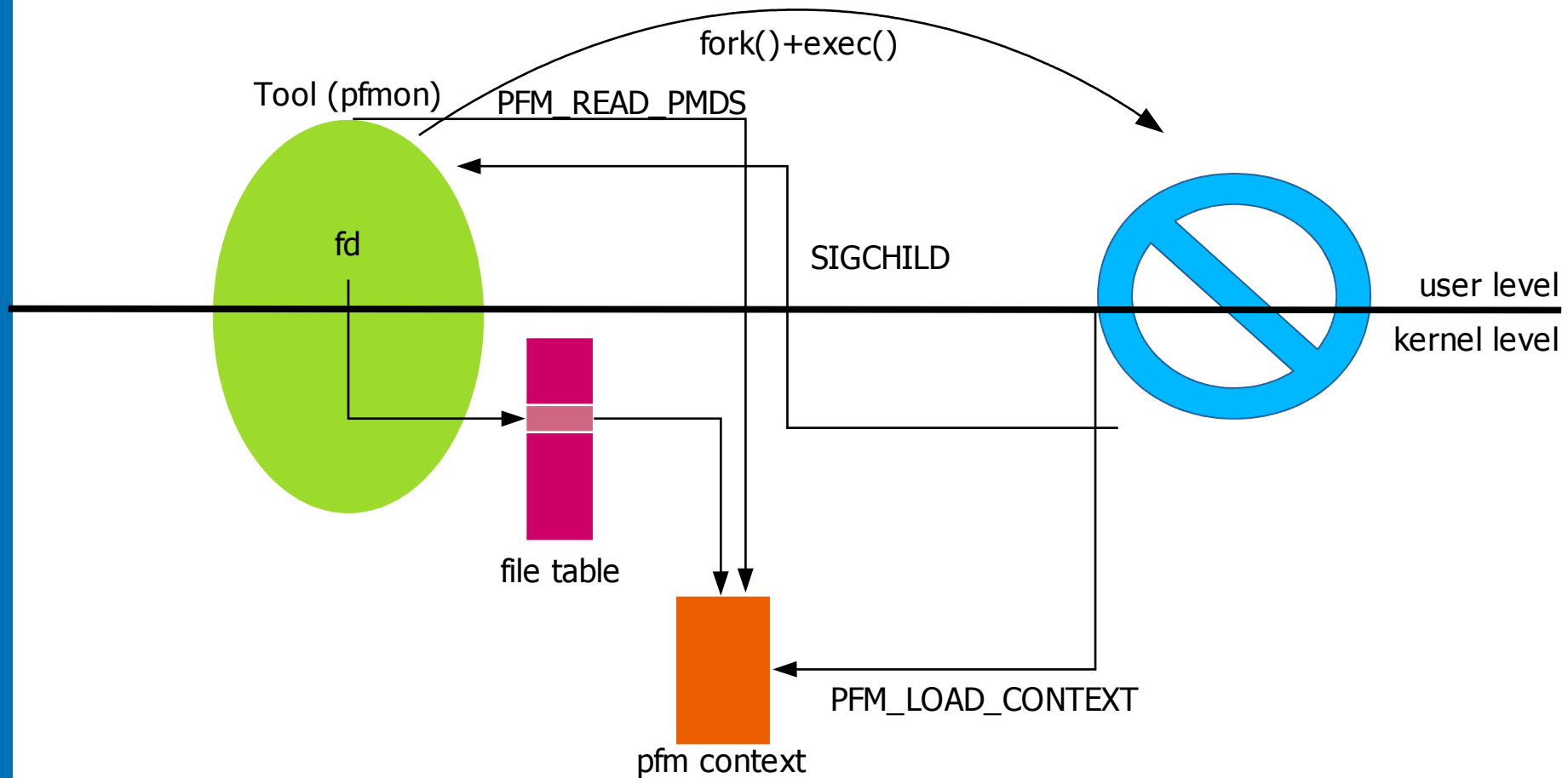
```
pfarg_context_t ctx;
pfarg_reg_t pmcs[1], pmds[1];
pfarg_load_t load_args;
pfmlib_input_param_t inp;
pfmlib_output_param_t outp;
pfm_find_event("CPU_CYCLES", &inp.pfp_events[0]);
inp.pfp_plm = PFM_PLM3; inp.pfp_count = 1;
pfm_dispatch_events(&inp, NULL, &outp, NULL);
pmcs[0].reg_num = pmcs[0].reg_num = outp.pfp_pmcs[0].reg_num;
pmcs[0].reg_value = outp.pfp_pmcs[0].reg_value;

perfmonctl(0, PFM_CREATE_CONTEXT, &ctx, 1);
fd = ctx.ctx_fd;
perfmonctl(fd, PFM_WRITE_PMCS, pmcs, 1);
perfmonctl(fd, PFM_WRITE_PMDs, pmds, 1);
load_arg.load_pid = getpid();
perfmonctl(fd, PFM_LOAD_CONTEXT, &load_args, 1);
perfmonctl(fd, PFM_START, NULL, 0);

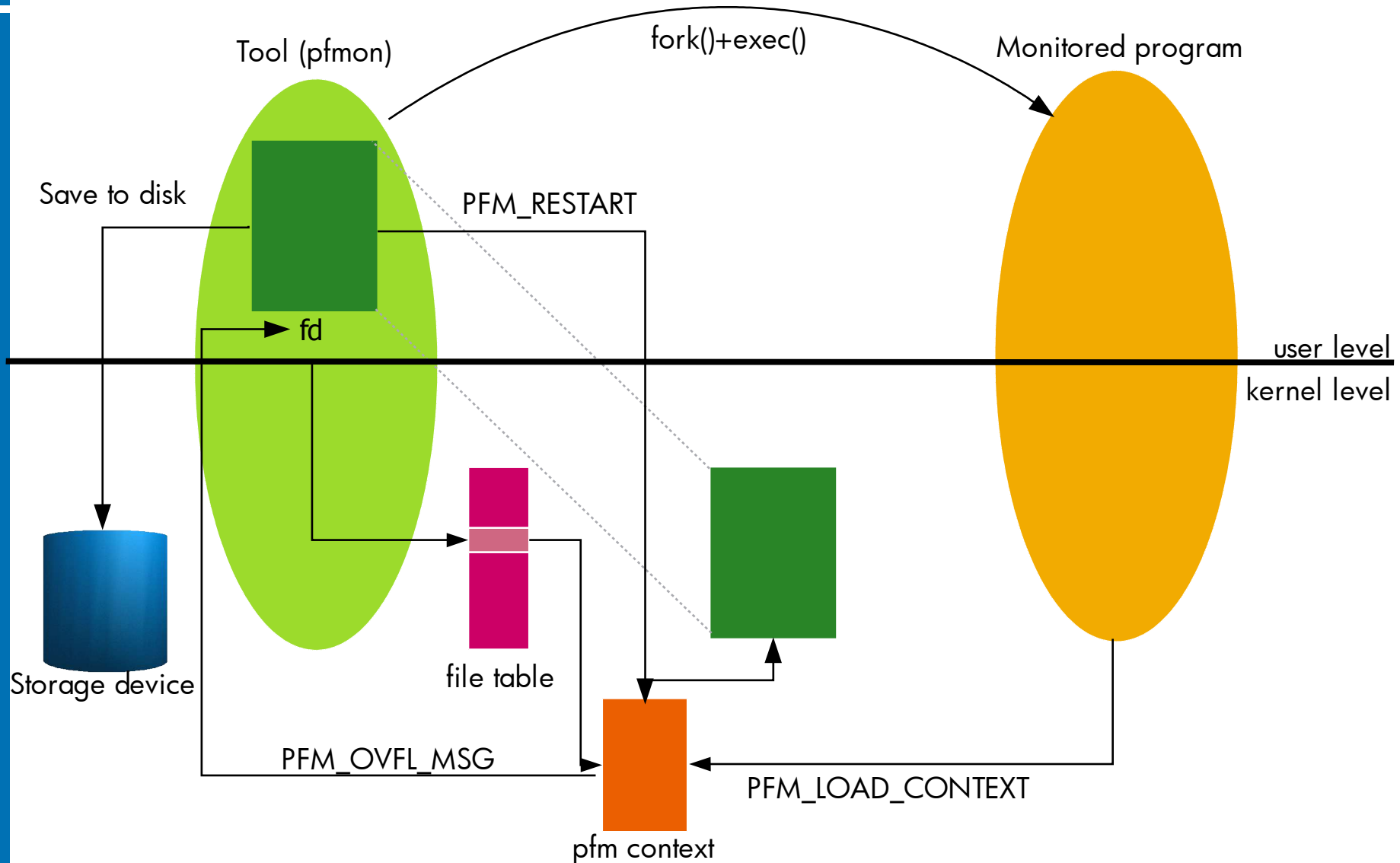
/* run code to measure */

perfmonctl(fd, PFM_STOP, NULL, 0);
perfmonctl(fd, PFM_READ_PMDs, pmds, 1);
close(fd);
```

Monitoring an unmodified binary



Typical sampling setup



Event-based sampling with pfmon

- Can sample on any event
- Multiple simultaneous sampling periods supported
 - Can collect several profiles at the same time
- Sampling period can be randomized
 - Very important to avoid biased results
- Ex.: every 100,000 cycles, record ia64_inst_retired

```
$ pfmon -reset-non-smpl -long-smpl-periods=100000 -ecpu_cycle,  
ia64_inst_retired -- foo
```

```
entry 6138 PID:4220 CPU:0 STAMP:0x480af599a15a IIP:0x400000000000003a0  
OVFL: 4 LAST_VAL: 100000  
PMD5 : 0x00000000000036ec4 (224964 instructions retired)  
entry 6139 PID:4220 CPU:0 STAMP:0x480af59b2c83 IIP:0x400000000000003a0  
OVFL: 4 LAST_VAL: 100000  
PMD5 : 0x00000000000036ec4
```

Sampling cache and TLB misses (EARS)



- Very useful to find where cache/TLB misses occur
 - Cannot be done with naïve IP-based sampling
 - Pinpoint the source of a miss, not the consequence
 - Must be used with sampling
- Ex.: sample every 1000 cache misses with latency > 4 cycles

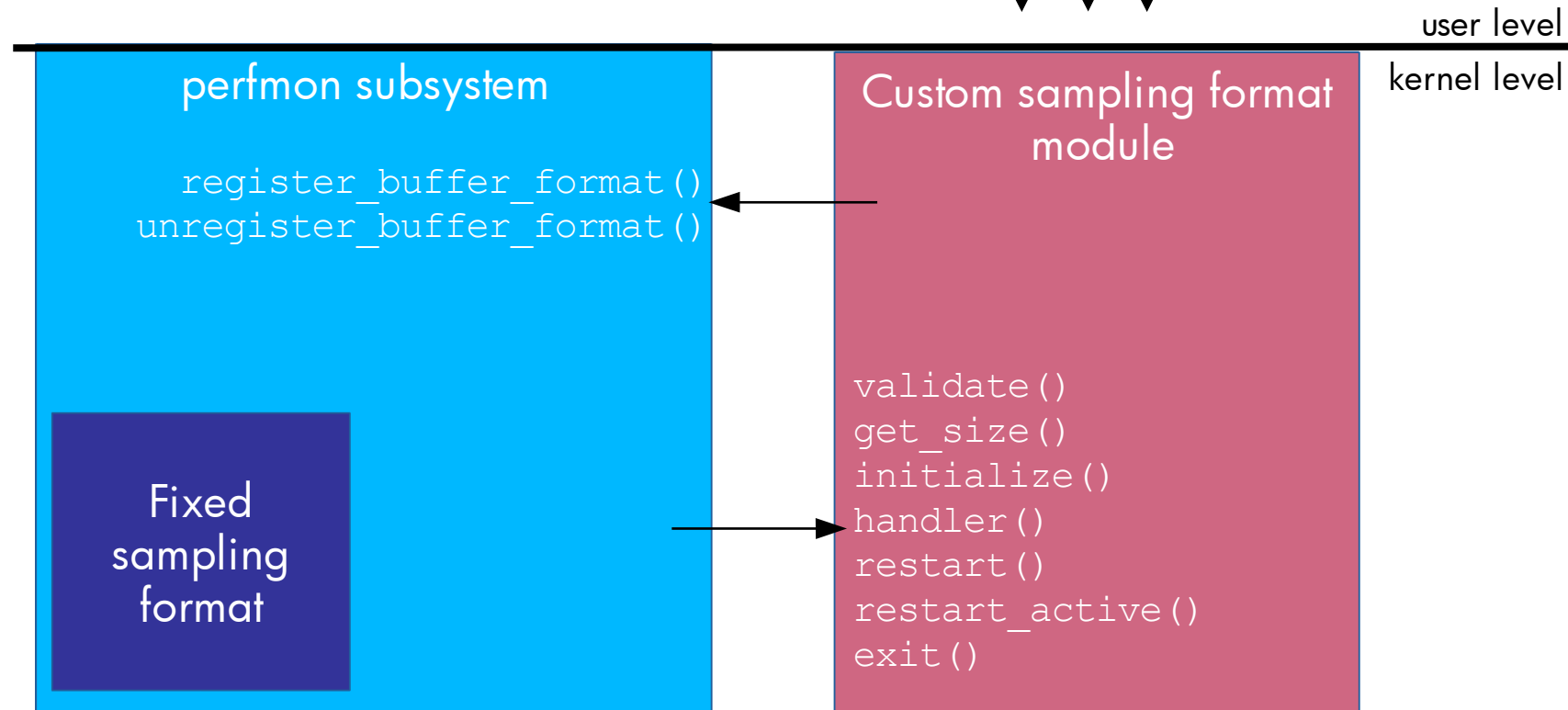
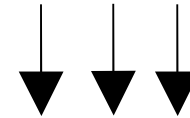
```
$ pfmon --long-smpl-periods=1000 -edata_ear_cache_lat4 - foo
```

```
entry 2000 PID:608 CPU:0 STAMP:0xfe3e1212e5 IIP:0x40000000000000990
    accessed data: 0x20000000000357000
    miss latency : 16 cycles
    inst address : 0x40000000000000981
```

```
40000000000000980:      [MMI]      ld8 r15=[r16]
0x40000000000000981:      ld8 r14=[r17] ← miss source
40000000000000982:      nop.i 0x0;;
40000000000000990:      [MMI]      cmp.ltu p7,p6=r14,r15;; ← stall
```

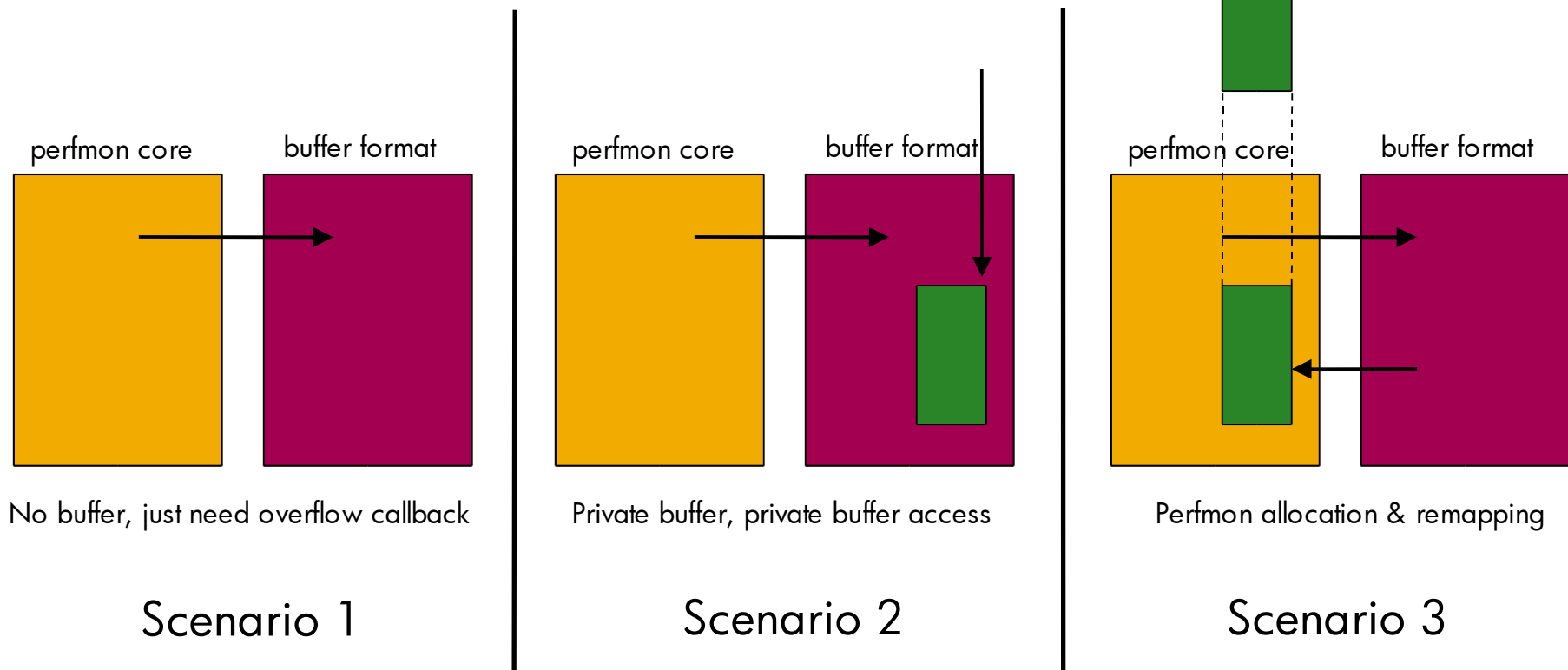
Custom sampling buffer interface

Format private interface
(optional)



Sampling buffer formats scenarios

- Formats can manage their own private buffers
- Formats can export their data anyway they want
- OProfile ported using scenario 2



pfmon/libpfm for perfmon-2

- Pfmon-3.0 (GNU GPL):
 - Monitoring of unmodified binaries or system-wide
 - Supports counting and event-based sampling
 - Supports all PMU features of Itanium® 1 & 2
 - Uses libpfm for setup
- Libpfm-3.0 (MIT-style):
 - Helps setup PMC registers
 - Provides common interface across PMU models
 - Contains all PMU specific knowledge (event encoding,...)
 - Provides model specific interface for advanced features
 - Does not invoke perfmonctl()
- Use pfmon-2.0/libpfm-2.0 for all 2.4 kernels



i n v e n t